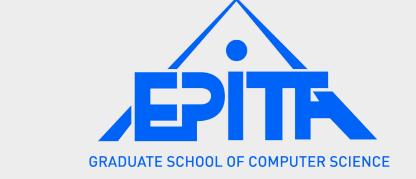


A 1ST PARALLEL ALGORITHM TO COMPUTE THE MORPHOLOGICAL TREE OF SHAPES OF $n\mathsf{D}$ IMAGES

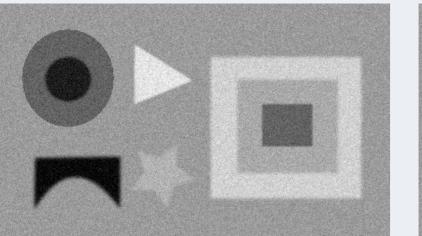
Sébastien Crozet Thierry Géraud*

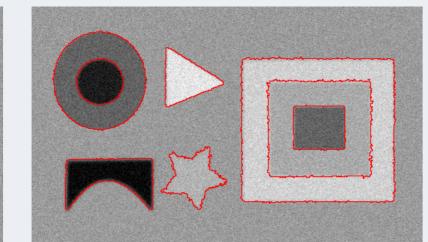


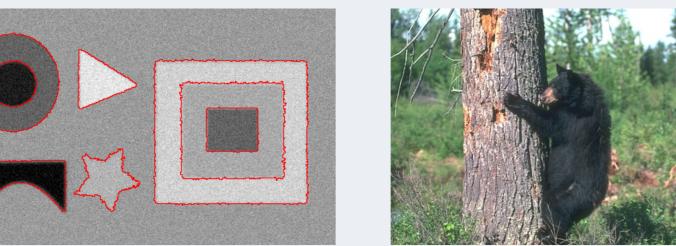
EPITA Research and Development Laboratory (LRDE), France thierry.geraud@lrde.epita.fr

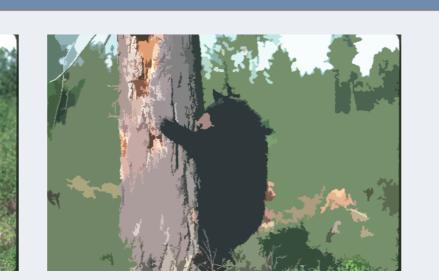
* also with Université Paris-Est, Laboratoire d'Informatique Gaspard-Monge (LIGM), Équipe A3SI, ESIEE Paris, France

The Tree of Shapes [1,4] as a Versatile Tool

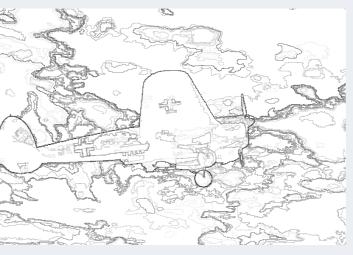


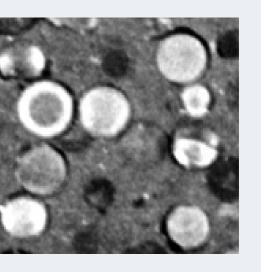


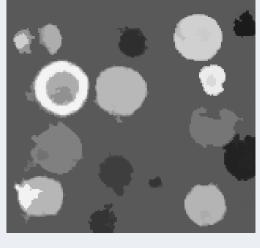












Object Detection (ICIP 2012)

Energy-Driven Simplification (ICIP 2013)

Hierarchy of Segmentations (ISMM 2013)

Shape Filtering (ICPR 2012)

All those results are from Yongchao Xu: http://www.lrde.epita.fr/wiki/User:Xu

At a Glance

Problem statement

- tree of shapes = self-dual morphological tree-based image representation
- a quasi-linear algorithm exists [5], yet it is sequential

Why is it interesting

- tree = easy structure to deal with
- nice properties: invariance to contrast changes and inversion
- numerous and powerful applications (see the banner above)

What our solution achieves

- ullet a 1st parallel version of the quasi-linear algorithm, and ready for $n{\sf D}$
- increasing size of data to process → no problemo :)

What follows from our solution

soon, processing 3D images with powerful self-dual morphological tools...

Algorithmic Scheme of the Sequential Version [5]

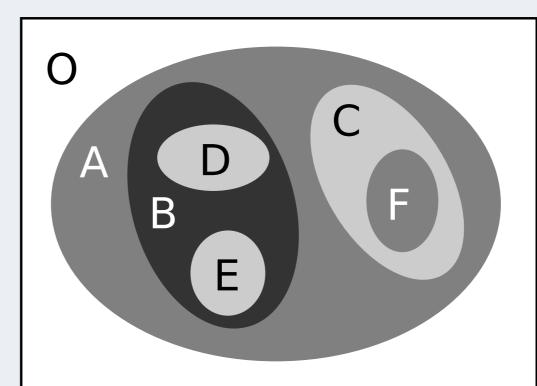
```
function ComputeTree(f, p_{\infty})
     \mathcal{F} \leftarrow \text{IMMERSE}(f)
      (\mathcal{R}, \mathcal{F}^{\flat}) \leftarrow \text{Sort}(\mathcal{F}, p_{\infty})
     par \leftarrow UnionFind(reverse(\mathcal{R}))
     return Canonicalize(par, \mathcal{R}, \mathcal{F}^{\flat})
end function
```

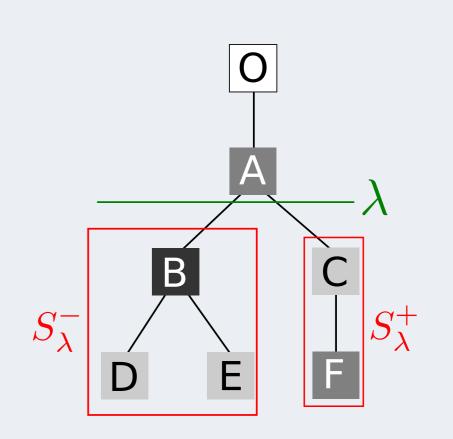
Algorithmic Scheme of the Parallel Version

```
function ComputeTree(f, p_{\infty})
     \mathcal{F} \leftarrow \text{ParallelImmerse}(f)
                                                         trivial
     \lambda \leftarrow \text{mean}(\mathcal{F}(p_{\infty}))
     Q[\lambda] \leftarrow p_{\infty}
     \mathcal{F}^{\text{ord}} \leftarrow \text{ParallelSort}(\mathcal{F}, Q, \lambda, 0)
                                                                    ▶ see [2] and [3]
     par \leftarrow ParallelMaxTree(\mathcal{F}^{ord})
     return Canonicalize(par, \mathcal{F}^{\mathrm{ord}})
end function
```

Parallel Sort NEW!

```
procedure ParallelSort(\mathcal{F}, Q, \mathcal{F}^{\mathrm{ord}}, \lambda, ord)
     Q[\lambda] \leftarrow p_{\infty}
     while any queue of Q is not empty do
          while Q[\lambda] is not empty do
               p \leftarrow \text{Pop}(Q[\lambda]), \quad \mathcal{F}^{\text{ord}}(p) \leftarrow ord
                for all n \in \mathcal{N}_4(p) that has not been visited yet do
                     if \lambda \in \mathcal{F}(n) then \text{Push}(Q[\lambda], n)
                     else if \lambda < \min(\mathcal{F}(n)) then \text{Push}(Q[\min(\mathcal{F}(n))], n)
                     else Push(Q[max(\mathcal{F}(n))], n)
                     end if
                end for
          end while
           ord \leftarrow ord + 1
          S_{\lambda}^{-} \leftarrow Q[0..\lambda], \quad S_{\lambda}^{+} \leftarrow Q[\lambda...\text{max value}]
          \lambda' \leftarrow highest level having faces on S_{\lambda}^{-}
           Run ParallelSort(\mathcal{F}, S_{\lambda}^{-}, \mathcal{F}^{\mathrm{ord}}, \lambda', \mathit{ord}) on another thread.
                                                                      \triangleright This thread continues with S_{\lambda}^{+}
          Q \leftarrow S_{\lambda}^{+}
          \lambda \leftarrow smallest level having faces on S_{\lambda}^{+}
     end while
     Wait for all child processes.
end procedure
                                                   example
```





An image and its tree of shapes. The nodes $oldsymbol{O}$ and $oldsymbol{A}$ have already been visited. The hierarchical queue contains the interior contour of $oldsymbol{B}$ and $oldsymbol{C}$. It is partitioned in two sets $S_{\lambda}^{+} = \partial B$ and $S_{\lambda}^{-} = \partial C$.

Reproducible Research

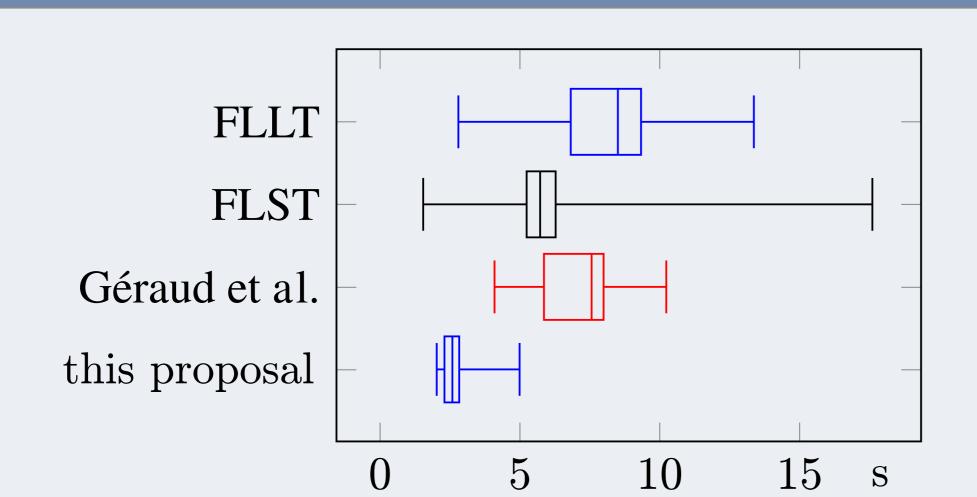
Evangelization from the Church of Mathematical Morphology

```
our C++ image processing library "Milena"
full source code of our method
```

- → http://olena.lrde.epita.fr
- → http://publis.lrde.epita.fr/crozet.14.icip →



Comparison



Computation times (in seconds) on a classical image test set of the following algorithms: FLLT [1], FLST [4], Géraud et al. [5], and this paper proposal.

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