

Two-Automaton Accepting Run Search in Spot

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Two-Automaton Accepting Run Search in Spot

- 1 Introduction
 - The model checking toolchain
 - The Spot library
 - The problem
- 2 The Accepting Run Search
 - What we need to build
 - Example
- 3 Results
 - Speed up
 - Comparison with other implementations
- 4 Conclusion

Introduction

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Introduction

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The model checking toolchain

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The Model Checking Toolchain

$$M \models \varphi$$

The Model Checking Toolchain

$$\begin{aligned} M \models \varphi \\ \iff L(A_M) \subseteq L(A_\varphi) \end{aligned}$$

The Model Checking Toolchain

$$\begin{aligned} & M \models \varphi \\ \iff & L(A_M) \subseteq L(A_\varphi) \\ \iff & L(A_M \otimes \overline{A_\varphi}) = \emptyset \end{aligned}$$

The Model Checking Toolchain

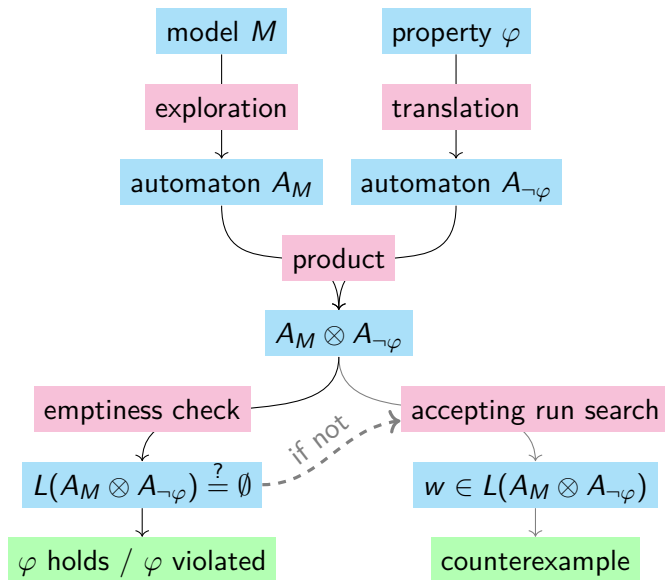
$$M \models \varphi$$

$$\iff L(A_M) \subseteq L(A_\varphi)$$

$$\iff L(A_M \otimes \overline{A_\varphi}) = \emptyset$$

$$\iff L(A_M \otimes A_{\neg\varphi}) = \emptyset$$

The Model Checking Toolchain



The Spot library

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- **The Spot library**
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- “Spot 2.0 — a framework for LTL and ω -automata manipulation” [4]
- C++ library
- algorithms, bridges for model checking
- executables for testing and comparing
 - `ltlcross`, `autcross`:
language equivalence



The problem

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The Problem

```
info: check_empty P1*N0
info: check_empty P1*N1
info: building Comp(N1)*Comp(P1) requires more acceptance sets than supported
info: building state-space #0/1 of 200 states with seed 0
info: state-space has 4225 edges
```

- ω -automata in Spot cannot have more than fixed number of acceptance sets (32 by default)
- acceptance sets sizes add-up during product

Operands may be expressed but not the product

Last Year's Work

Two-automaton emptiness check:

- Two-automaton emptiness check in Spot [5]
- Simulate a product on-the-fly
- Based on Couvreur's on-the-fly emptiness check algorithm [2]
- Optimisations on automata strength and API

No accepting run search, we cannot use it to look for a counterexample.

The Accepting Run Search

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What we need to build

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What We Need To Build

Pumping lemma: infinite word, but finite automata

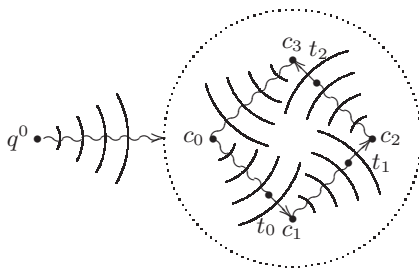


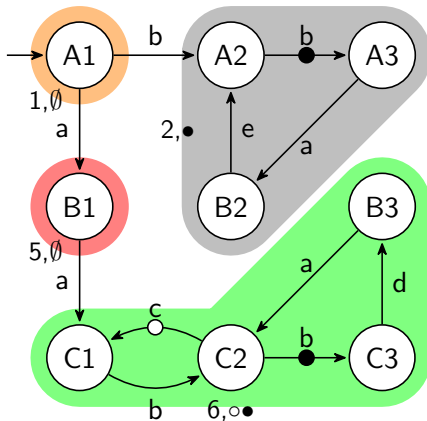
Fig. 5. Computing an accepting run for a TGBA.

Excerpt from “On-the-Fly Emptiness Checks for Generalized Büchi Automata” [3]

Example

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Example



$$Inf(\bullet) \wedge Inf(\circ)$$

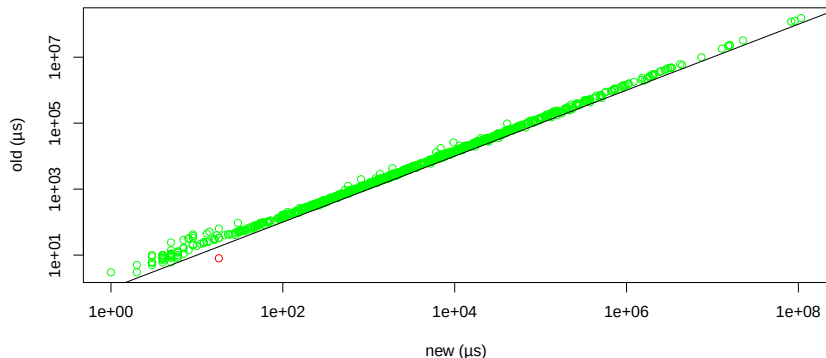
Results

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Speed up

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Comparison of old vs. new emptiness check in ltlcross



- 100 formulae¹ × 3 tools² ⇒ 849 emptiness checks (all empty, no accepting run search) + 27 timeouts
- outlier: "`(!(X(p11))) -> (!(p0))`" through 1t13ba

¹randtl --tree-size=30 30

`21tlcross -T 60 1tl2ba 1tl3ba 1tl2tgba`

Comparison with other implementations

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- Speed up

- Comparison with other implementations

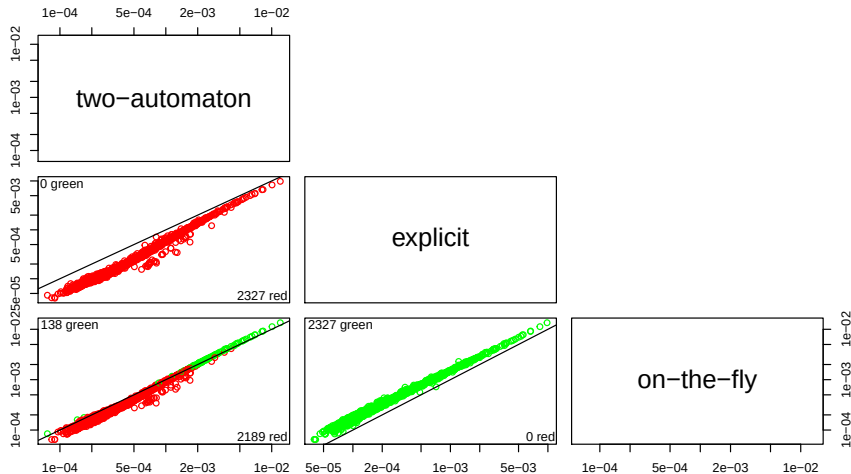
4 Conclusion

Bench Setup

- 1 We measure the time spent in:
 - two-automaton emptiness check
 - explicit product + emptiness check
 - on-the-fly product + emptiness check
 - + accepting run search when not empty
- 2 Data set:
 - 100 random automata: 200 states, 16 acceptance sets, 10 condition variables
 - 5050 combinations, 2668 empty, 2382 non-empty
- 3 In parallel over 24 cores of the same machine

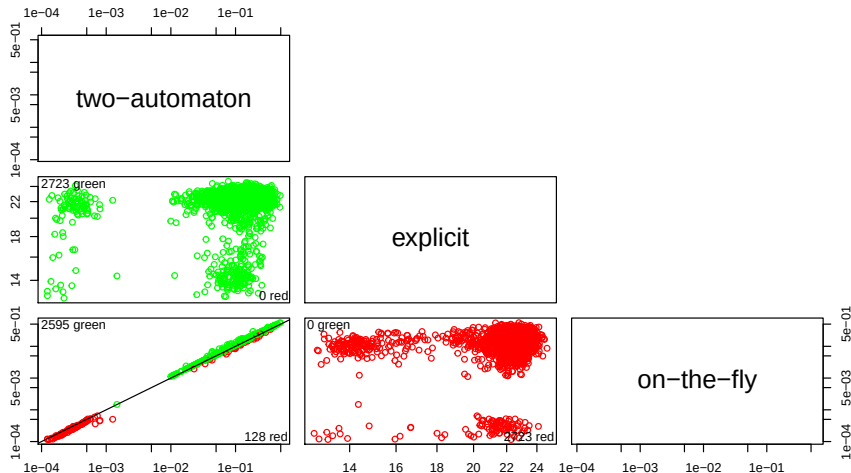
Case of Emptiness

Time of positive emptiness check (s)



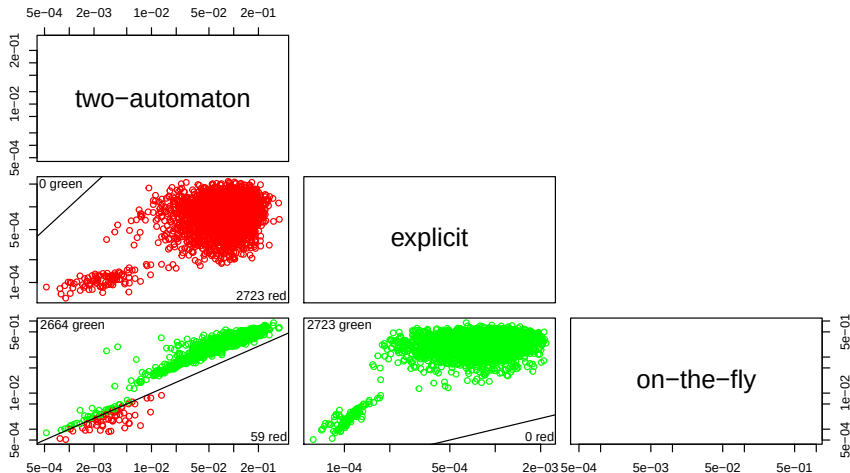
Case of Non-Emptiness

Time of negative emptiness check (s)



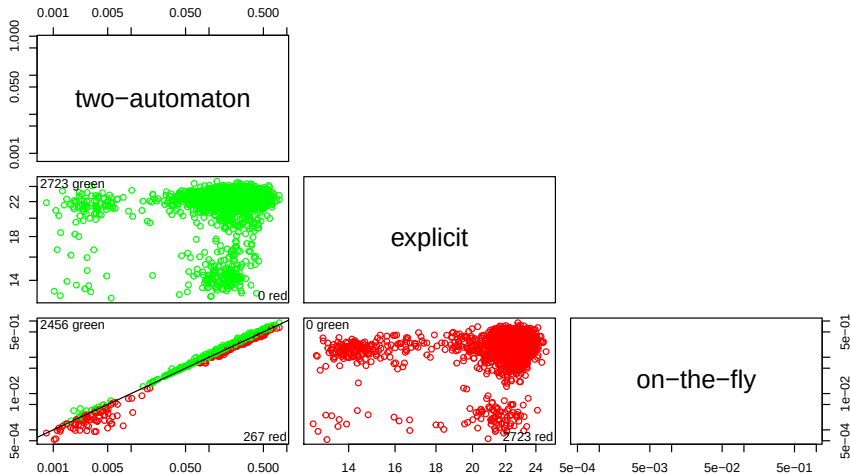
Case of Non-Emptiness

Time of accepting run search (s)



Case of Non-Emptiness

Time of negative emptiness check + accepting run search (s)



Conclusion

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Conclusion

What has been done:

- Make the two-automaton emptiness check fully usable
- Test it on a broader range

What remains to be done:

- Investigate on performances
- From generalised Büchi to Rabin: “Explicit State Model Checking with Generalized Büchi and Rabin Automata” [1]

Questions?

Bibliography I



Vincent Bloemen, Alexandre Duret-Lutz, and Jaco van de Pol. “Explicit State Model Checking with Generalized Büchi and Rabin Automata”. In: [Proceedings of the 24th International SPIN Symposium on Model Checking of Software \(SPIN’17\)](#). ACM, July 2017, pp. 50–59.



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Bibliography III



Clément Gillard. [Two-automaton emptiness check in Spot](#). Tech. rep. 1706. EPITA Research and Development Laboratory (LRDE), 2017.