# **Compiler Construction**

 $\sim$  Syntactic Analysis  $\checkmark$ 

# **Syntactic Analysis**

Reminders from the THL lectures!

How to check that a text (stream of token) is valid according to a given grammar?  $\Rightarrow$  We need something more powerful than finite automata!



#### Language

#### A **Language** is a set of strings, each string is a sequence of **symbols** from an **alphabet**

#### Remark

In our context, the alphabet is the set of token types returned by the lexical analyzer.

# Context-free grammars (1/2)

**Regular expressions** 

Regular expressions are not enough to represent programming languages.

#### Context-free Grammar

A context-free grammar is a set of recursive rules used to generate patterns of symbols.

# **Context-free grammars (2/2)**

#### **Production rules**

A grammar has production rules of the form symbol  $\rightarrow$  symbol symbol symbol Symbols can be:

- terminal, meaning that this is a token from the alphabet.
- non-terminal, meaning that it appears on the left-hand side of some production.

Example	
E -> E + E	
E -> id	

### **Derivations**

To detect if a sentence is in the expected language we can perform **derivations**:

- Start with a symbol
- Apply productions rules (Replace any non terminal by one of its right-hand side)
- Repeat until no more replacement

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There are many derivations:

- the leftmost non-terminal symbol is always expanded
- the **rightmost** non-terminal symbol is always expanded

#### **Parse tree**

#### Parse tree

A **parse tree** is made by connecting each symbol in a derivation to the one from which it was derived

Example input: id + id + id(One possible) derivation:  $\underline{\underline{E}}$  $\underline{\underline{E}} + \underline{\underline{E}}$  $\underline{\underline{E}} + id$  $\underline{\underline{E}} + id$  $\underline{\underline{E}} + id + id$ id + id + id



# **Ambiguous Grammar**

#### Ambiguous Grammar

A grammar is **ambiguous** if we can derive two different parse tree for a sentence



### **Predictive Parsing – Unambiguous Grammar**

A predictive parser is a recursive descent parser which is able to predict which production rule is used to replace the input string.  $\Rightarrow$  each sub-expression must provide enough information to choose a production rule.

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- LL(k): Left-to-right, Leftmost derivation
- LR(k): Left-to-right, Rightmost derivation

**Summary** 

